On the data processing inequality for the relative entropy between two quantum states

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Joint work with Andreas Bluhm (U. Copenhagen)

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Origins of Quantum Information Theory







Shannon



Bennett









$$H |\psi\rangle = E |\psi\rangle$$
$$\Delta x \Delta p \ge \frac{\hbar}{2}$$

"Nature isn't classical, dammit, and if you want to make a simulation of nature. you'd better make it quantum mechanical"



Given n = pq. Find (p, q)



Quantum supremacy



New algorithms



1920s

First quantum revolution

New technologies

New applications

The Bell System Technical Journal

Vol. XXVII July, 1945 A Mathematical Theory of Communication By C. E. SHANNON

$$H(p) := -\sum_{x} p(x) \log p(x)$$

1948

1984 1994

Second quantum revolution

Classical digital revolution

Acknowledgement: David Sutter (IBM Research)

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$$I(X:Z|Y)_P \geq 0$$
 (trivial).



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Operational meaning of $D(\sigma||\rho) - D(\mathcal{T}(\sigma)||\mathcal{T}(\rho))$

- Thermodynamics: Cost of a certain quantum process (Faist et al, '18).
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Problem

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More specifically, if we consider $\mathcal{V}_{BC} \circ \mathcal{R}_{\text{tr}_C}^{\sigma_{BC}} \circ \mathcal{U}_B$, with U_B and V_{BC} unitaries on \mathcal{H}_B , \mathcal{H}_{BC} respectively,

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Extensions and improvements of the previous result:

$$D(\sigma||\rho) - D(\mathcal{T}(\sigma)||\mathcal{T}(\rho)) \ge (1), (2), (3), \text{ where:}$$

$$(1) := -\int \beta_0(t) \log F\left(\sigma, \mathcal{R}^{\rho, [t]}_{\mathcal{T}} \circ \mathcal{T}(\sigma)\right) dt \text{ (Junge et al. '15)},$$

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Answer: It is not possible (Brandao et al. '15, Fawzi² '17).

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$$D(\sigma_{ABC}||\mathcal{R}_{\operatorname{tr}_{C}}^{\sigma_{BC}} \circ \operatorname{tr}_{C}[\sigma_{ABC}]) + \Lambda_{\max}(\sigma_{AB}||\mathcal{R}_{B \to B}) \ge I(A:C|B)_{\sigma},$$

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$$\Lambda_{\max}(\sigma||\mathcal{E}) = 0 \Leftrightarrow \mathcal{E}(\sigma) = \sigma,$$

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PROBLEM

Can we find a lower bound for the DPI in terms of $D(\sigma||\mathcal{R}^{\rho}_{\tau} \circ \mathcal{T}(\sigma))$?

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Problem

Can we find a lower bound for the DPI in terms of $\mathcal{R}^{\rho}_{\tau} \circ \mathcal{T}(\sigma)$?

(Carlen-Vershynina '17) $\mathcal{E}: \mathcal{M} \to \mathcal{N}$ conditional expectation, $\sigma_{\mathcal{N}} := \mathcal{E}(\sigma)$ and $\rho_{\mathcal{N}} := \mathcal{E}(\rho)$:

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Some definitions

CONDITIONAL EXPECTATION

Let \mathcal{M} matrix algebra with matrix subalgebra \mathcal{N} . There exists a unique linear mapping $\mathcal{E}: \mathcal{M} \to \mathcal{N}$ such that

- \bullet \mathcal{E} is a positive map,
- $\mathcal{E}(B) = B \text{ for all } B \in \mathcal{N},$
- **3** $\mathcal{E}(AB) = \mathcal{E}(A)B$ for all $A \in \mathcal{M}$ and all $B \in \mathcal{N}$,
- **4** \mathcal{E} is trace preserving.

A map fulfilling (1)-(3) is called a *conditional expectation*.

Belavkin-Staszewski relative entropy

Given $\sigma > 0, \rho > 0$ states on a matrix algebra \mathcal{M} , their **BS-entropy** is defined as:

$$\hat{S}_{\mathrm{BS}}(\sigma||\rho) := \mathrm{tr} \Big[\sigma \log \Big(\sigma^{1/2} \rho^{-1} \sigma^{1/2} \Big) \Big].$$

RELATION BETWEEN BELATIVE ENTROPIES

The following holds for every $\sigma > 0, \rho > 0$

$$\hat{S}_{\mathrm{BS}}(\sigma||\rho) \ge D(\sigma||\rho)$$

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Some definitions

OPERATOR CONVEX

Let $\mathcal{I} \subseteq \mathbb{R}$ interval and $f: \mathcal{I} \to \mathbb{R}$. If

$$f(\lambda A + (1 - \lambda)B) \le \lambda f(A) + (1 - \lambda)f(B)$$

for all Hermitian $A, B \in \mathcal{B}(\mathcal{H})$ with spectrum contained in \mathcal{I} , all $\lambda \in [0, 1]$, and for all finite-dimensional Hilbert spaces \mathcal{H} , then f is operator convex.

(Hiai-Mosonyi '17)

STANDARD f-DIVERGENCES

Let $f:(0,\infty)\to\mathbb{R}$ be an operator convex function and $\sigma>0,\ \rho>0$ be two states on a matrix algebra \mathcal{M} . Then,

$$S_f(\sigma||\rho) = \operatorname{tr}\left[\rho^{1/2} f(L_{\sigma} R_{\rho^{-1}}) \rho^{1/2}\right]$$

is the standard f-divergence.

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Example: Let $f(x) = x \log x$. Then,

$$S_f(\sigma||\rho) = \operatorname{tr}[\sigma(\log \sigma - \log \rho)]$$

defines the relative entropy $D(\sigma || \rho)$.

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$$S_f(\mathcal{T}(\sigma)||\mathcal{T}((\rho)) \le S_f(\sigma||\rho)$$

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CONDITIONS FOR EQUALITY

Let $\sigma > 0$, $\rho > 0$ be on \mathcal{M} and let $\mathcal{T} : \mathcal{M} \to \mathcal{B}$ be a 2PTP linear map. Then, the following are equivalent:

- There exists a TP map $\hat{\mathcal{T}}: \mathcal{B} \to \mathcal{M}$ such that $\hat{\mathcal{T}}(\mathcal{T}(\rho)) = \rho$ and $\hat{\mathcal{T}}(\mathcal{T}(\sigma)) = \sigma$.
- $S_f(\mathcal{T}(\sigma)||\mathcal{T}(\rho)) = S_f(\sigma||\rho)$ for all operator convex f on $[0,\infty)$.

Maximal f-divergences

Let $f:(0,\infty)\to\mathbb{R}$ be an operator convex function and $\sigma>0,\,\rho>0$ be two states on a matrix algebra \mathcal{M} . Then,

$$\hat{S}_f(\sigma \| \rho) = \text{tr}\left[\rho^{1/2} f(\rho^{-1/2} \sigma \rho^{-1/2}) \rho^{1/2}\right]$$

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Let $\sigma > 0$, $\rho > 0$ be on \mathcal{M} and $\mathcal{T} : \mathcal{M} \to \mathcal{B}$ be a PTP linear map. Then, the following are equivalent:

- $\hat{S}_f(\mathcal{T}(\sigma) \| \mathcal{T}(\rho)) = \hat{S}_f(\sigma \| \rho) \text{ for all operator convex functions on } [0, \infty).$
- $\mathbf{2} \operatorname{tr} \left[\mathcal{T}(\sigma)^2 \mathcal{T}(\rho)^{-1} \right] = \operatorname{tr} \left[\sigma^2 \rho^{-1} \right].$

Relation between f-divergences

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For every two states $\sigma > 0, \, \rho > 0$ on \mathcal{M} and every operator convex function $f:(0,\infty) \to \mathbb{R}$,

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REMARK: DIFFERENCE

For maximal f-divergences, there is no equivalent condition for equality in DPI which provides a explicit expression of recovery for σ .

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QUESTIONS

BS RECOVERY CONDITION

Can we prove an equivalent condition for equality in DPI for the BS entropy (or for maximal f-divergences) which provides a explicit expression of recovery for σ ?

STRENGTHENED DPI FOR BS ENTROPY

Following Carlen-Vershynina, can we provide a lower bound for the DPI for the BS entropy (or for maximal f-divergences) in terms of a (hypothetical) BS recovery condition?

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Equivalent conditions for equality on DPI

$$\Gamma := \sigma^{-1/2} \rho \sigma^{-1/2} \text{ and } \Gamma_{\mathcal{T}} := \sigma_{\mathcal{T}}^{-1/2} \rho_{\mathcal{T}} \sigma_{\mathcal{T}}^{-1/2}$$
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Equivalent conditions for equality on DPI (Bluhm-C. '19

Let \mathcal{M} be a matrix algebra with unital subalgebra \mathcal{N} . Let $\mathcal{T}: \mathcal{M} \to \mathcal{N}$ be the trace-preserving conditional expectation onto this subalgebra. Let $\sigma > 0$, $\rho > 0$ be two quantum states on \mathcal{M} . Then, the following are equivalent:

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Consequences

Note: Although they can be seen as a consequence of the previous result, the following facts were previously known.

COROLLARY

$$\hat{S}_{BS}(\sigma \| \rho) = \hat{S}_{BS}(\sigma_{\mathcal{T}} \| \rho_{\mathcal{T}}) \Leftrightarrow \rho = \mathcal{B}_{\mathcal{T}}^{\sigma} \circ \mathcal{T}(\rho)$$

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Corollary

$$D(\sigma \| \rho) = D(\sigma_{\mathcal{T}} \| \rho_{\mathcal{T}}) \implies \hat{S}_{BS}(\sigma \| \rho) = \hat{S}_{BS}(\sigma_{\mathcal{T}} \| \rho_{\mathcal{T}}).$$

Equivalently

$$\sigma = \mathcal{R}^{\rho}_{\mathcal{T}} \circ \mathcal{T}(\sigma) \implies \sigma = \mathcal{B}^{\rho}_{\mathcal{T}} \circ \mathcal{T}(\sigma)$$

The converse of this result is false (Jencová-Petz-Pitrik '09, Hiai-Mosonyi '17).

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RESULTS FOR THE BS-ENTROPY, Bluhm-C. '19

Relative entropy	BS-entropy
$\operatorname{tr}[\sigma(\log\sigma-\log\rho)]$	$\operatorname{tr} \bigl[\sigma \log \left(\sigma^{1/2} \rho^{-1} \sigma^{1/2} \right) \bigr]$
$\rho = \rho^{1/2} \mathcal{T}^* \left(\mathcal{T}(\rho)^{-1/2} \mathcal{T}(\sigma) \mathcal{T}(\rho)^{-1/2} \right) \rho^{1/2}$	$\sigma = \rho \mathcal{T}^* \left(\mathcal{T}(\rho)^{-1} \mathcal{T}(\sigma) \right)$
$\left(\frac{\pi}{8}\right)^4 \ L_{\rho}R_{\sigma^{-1}}\ _{\infty}^{-2} \ \mathcal{R}_{\mathcal{E}}^{\sigma}(\rho_{\mathcal{N}}) - \rho\ _{1}^{4}$	$\left(\frac{\pi}{8}\right)^{4} \ \Gamma\ _{\infty}^{-4} \ \sigma^{-1}\ _{\infty}^{-2} \ \rho - \mathcal{B}_{\mathcal{T}}^{\sigma} \circ \mathcal{T}(\rho)\ _{2}^{4}$
Extension to standard f-divergences	Extension to maximal f-divergences



Particular case: $\mathcal{H}_{ABC} = \mathcal{H}_A \otimes \mathcal{H}_B \otimes \mathcal{H}_C$.

Quantum channel: $\mathcal{T} = \operatorname{tr}_{\mathcal{C}}$.



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$$\sigma = \mathcal{R}^{\rho}_{\mathcal{T}} \circ \mathcal{T}(\sigma) \leadsto \sigma_{ABC} = \sigma_{BC}^{1/2} \, \sigma_{B}^{-1/2} \, \sigma_{AB} \, \sigma_{B}^{-1/2} \, \sigma_{BC}^{1/2}.$$



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Define a BS quantum state as a state $\sigma_{ABC} \in \mathcal{S}_{ABC}$ such that $\sigma_{ABC} = \sigma_{BC} \sigma_B^{-1} \sigma_{AB}$.



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QUESTION

Is the set of BS quantum states robust?



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$$\sigma_{ABC} = \sigma_{BC}^{1/2} \sigma_{B}^{-1/2} \sigma_{AB} \sigma_{B}^{-1/2} \sigma_{BC}^{1/2}$$
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